CHAPTER 10 SSE4.2 AND SIMD PROGRAMMING FOR TEXT-PROCESSING/LEXING/PARSING

String/text processing spans a discipline that often employs techniques different from traditional SIMD integer vector processing. Much of the traditional string/text algorithms are character based, where characters may be represented by encodings (or code points) of fixed or variable byte sizes. Textual data represents a vast amount of raw data and often carrying contextual information. The contextual information embedded in raw textual data often requires algorithmic processing dealing with a wide range of attributes, such as character values, character positions, character encoding formats, subsetting of character sets, strings of explicit or implicit lengths, tokens, delimiters; contextual objects may be represented by sequential characters within a pre-defined character subsets (e.g. decimal-valued strings); textual streams may contain embedded state transitions separating objects of different contexts (e.g. tag-delimited fields).

Traditional Integer SIMD vector instructions may, in some simpler situations, be successful to speed up simple string processing functions. SSE4.2 includes four new instructions that offer advances to computational algorithms targeting string/text processing, lexing and parsing of either unstructured or structured textual data.

10.1 SSE4.2 STRING AND TEXT INSTRUCTIONS

SSE4.2 provides four instructions, PCMPESTRI/PCMPESTRM/PCMPISTRI/PCMPISTRM that can accelerate string and text processing by combining the efficiency of SIMD programming techniques and the lexical primitives that are embedded in these 4 instructions. Simple examples of these instructions include string length determination, direct string comparison, string case handling, delimiter/token processing, locating word boundaries, locating sub-string matches in large text blocks. Sophisticated application of SSE4.2 can accelerate XML parsing and Schema validation.

Processor's support for SSE4.2 is indicated by the feature flag value returned in ECX [bit 20] after executing CPUID instruction with EAX input value of 1 (i.e. SSE4.2 is supported if CPUID.01H:ECX.SSE4_2 [bit 20] = 1). Therefore, software must verify CPUID.01H:ECX.SSE4_2 [bit 20] is set before using these 4 instructions. (Verifying CPUID.01H:ECX.SSE4_2 = 1 is also required before using PCMPGTQ or CRC32. Verifying CPUID.01H:ECX.POPCNT[Bit 23] = 1 is required before using the POPCNT instruction.)

These string/text processing instructions work by performing up to 256 comparison operations on text fragments. Each text fragment can be 16 bytes. They can handle fragments of different formats: either byte or word elements. Each of these four instructions can be configured to perform four types of parallel comparison operation on two text fragments.

The aggregated intermediate result of a parallel comparison of two text fragments become a bit patterns:16 bits for processing byte elements or 8 bits for word elements. These instruction provide additional flexibility, using bit fields in the immediate operand of the instruction syntax, to configure an unary transformation (polarity) on the first intermediate result.

Lastly, the instruction's immediate operand offers a output selection control to further configure the flexibility of the final result produced by the instruction. The rich configurability of these instruction is summarized in Figure 10-1.

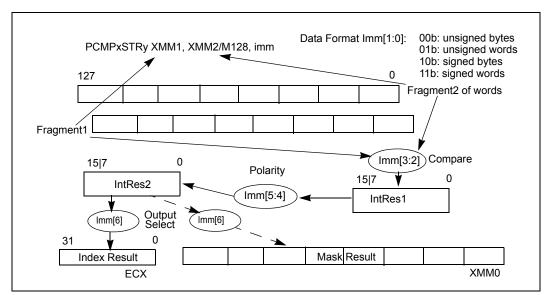


Figure 10-1. SSE4.2 String/Text Instruction Immediate Operand Control

The PCMPxSTRI instructions produce final result as an integer index in ECX, the PCMPxSTRM instructions produce final result as a bit mask in the XMM0 register. The PCMPISTRy instructions support processing string/text fragments using implicit length control via null termination for handling string/text of unknown size. the PCMPESTRy instructions support explicit length control via EDX:EAX register pair to specify the length text fragments in the source operands.

The first intermediate result, IntRes1, is an aggregated result of bit patterns from parallel comparison operations done on pairs of data elements from each text fragment, according to the imm[3:2] bit field encoding, see Table 10-1.

lmm[3:2]	Name	IntRes1[i] is TRUE if	Potential Usage	
00B	Equal Any	Element i in fragment2 matches any element j in fragment1	Tokenization, XML parser	
01B	Ranges	Element i in fragment2 is within any range pairs specified in fragment1	Subsetting, Case handling, XML parser, Schema validation	
10B	Equal Each	Element i in fragment2 matches element i in fragment1	Strcmp()	
11B	Equal Ordered	Element i and subsequent, consecutive valid elements in fragment2 match fully or partially with fragment1 starting from element 0		

Input data element format selection using imm[1:0] can support signed or unsigned byte/word elements.

The bit field imm[5:4] allows applying a unary transformation on IntRes1, see Table 10-2.

Imm[5:4]	Name	IntRes2[i] =	Potential Usage
00B	No Change	IntRes1[i]	
01B	Invert	-IntRes1[i]	
10B	No Change	IntRes1[i]	
11B	Mask Negative	IntRes1[i] if element i of fragment2 is invalid, otherwise - IntRes1[i]	

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Table 10-2. 3364.2 Str	ing/rext instructions unary	/ Transformation on IntRes1

The output selection field, imm[6] is described in Table 10-3.

lmm[6]	Instruction	Final Result	Potential Usage
0B	PCMPxSTRI	ECX = offset of least significant bit set in IntRes2 if IntRes2 != 0, otherwise	
		ECX = number of data element per 16 bytes	
0B	PCMPxSTRM	XMM0 = ZeroExtend(IntRes2);	
1B	PCMPxSTRI	ECX = offset of most significant bit set in IntRes2 if IntRes2 != 0, otherwise ECX = number of data element per 16 bytes	
1B	PCMPxSTRM	Data element i of XMM0 = SignExtend(IntRes2[i]);	

The comparison operation on each data element pair is defined in Table 10-4. Table 10-4 defines the type of comparison operation between valid data elements (last row of Table 10-4) and boundary conditions when the fragment in a source operand may contain invalid data elements (rows 1 through 3 of Table 10-4). Arithmetic comparison are performed only if both data elements are valid element in fragment1 and fragment2, as shown in row 4 of Table 10-4.

Table 10-4. SSE4.2 String/Text Instructions Element-Pair Comparison Definition

fragment1 element	fragment2 element	lmm[3:2]= 00B, Equal Any	lmm[3:2]= 01B, Ranges	lmm[3:2]= 10B, Equal Each	lmm[3:2]= 11B, Equal Ordered
invalid	invalid	Force False	Force False	Force True	Force True
invalid	valid	Force False	Force False	Force False	Force True
valid	invalid	Force False	Force False	Force False	Force False
valid	valid	Compare	Compare	Compare	Compare

The string and text processing instruction provides several aid to handle end-of-string situations, see Table 10-5. Additionally, the PCMPxSTRy instructions are designed to not require 16-byte alignment to simplify text processing requirements.

EFLAGs	Description	Potential Usage
CF	Reset if IntRes2 = 0; Otherwise set	When CF=0, ECX= #of data element to scan next
ZF	Reset if entire 16-byte fragment2 is valid	likely end-of-string
SF	Reset if entire 16-byte fragment1 is valid	
OF	IntRes2[0];	

Table 10-5. SSE4.2 String/Text Instructions Eflags Behavior